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(54) **SERIAL TURBO TRELLIS CODED  
MODULATION USING A SERIALY  
CONCATENATED CODER**

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13, 2000.

(51) **Int. Cl.**  
**H03M 13/03** (2006.01)  
**H03M 3/00** (2006.01)

(52) **U.S. Cl.** ..... **714/792; 714/755**

(58) **Field of Classification Search** ..... **714/792,**  
**714/755**

See application file for complete search history.

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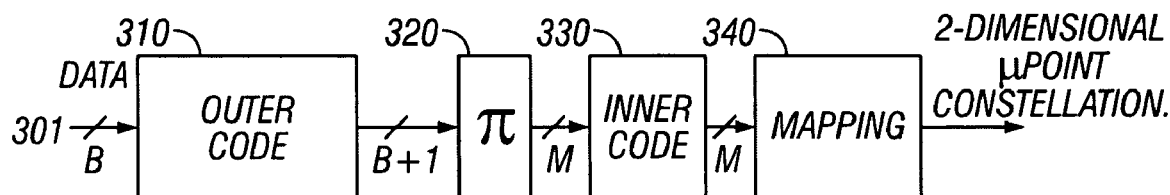
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*Primary Examiner*—Joseph D. Torres

(57) **ABSTRACT**

A coding system uses a serially concatenated coder driving  
an interleaver, which drives a trellis coder. This combina-  
tion, while similar to a turbo coder, produces certain differ-  
ent characteristics.

**16 Claims, 5 Drawing Sheets**



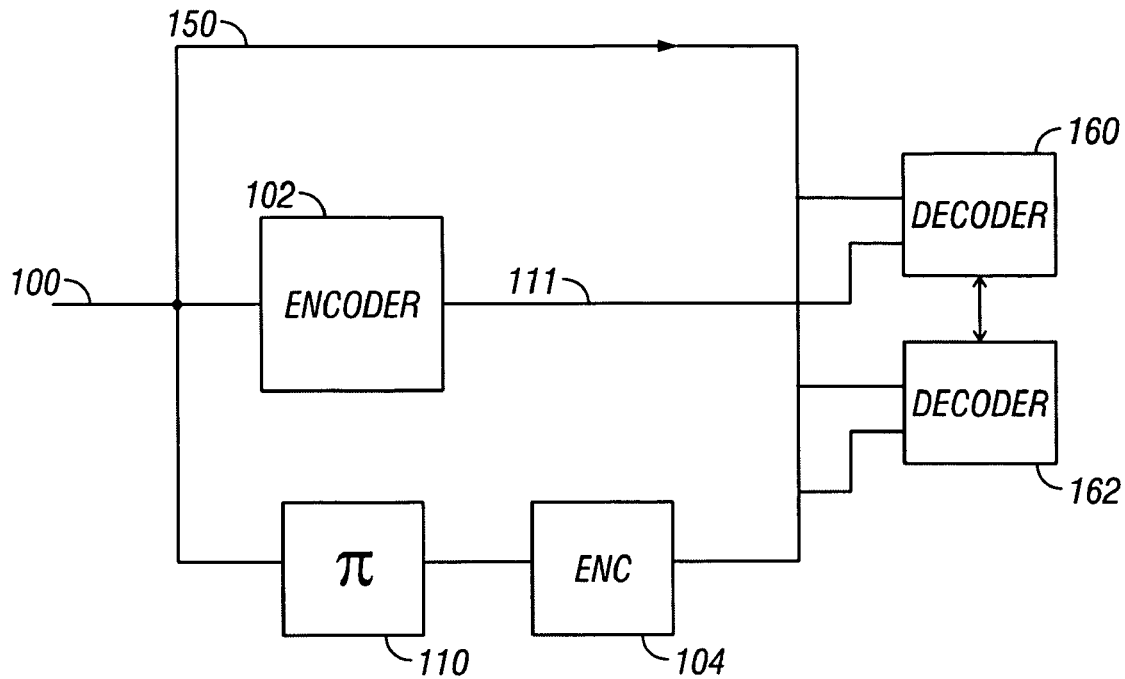


FIG. 1

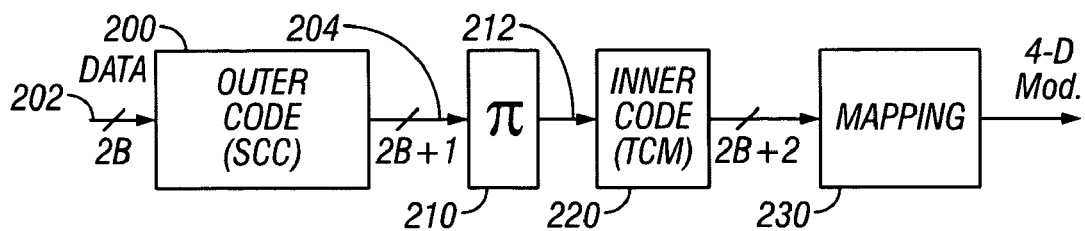


FIG. 2

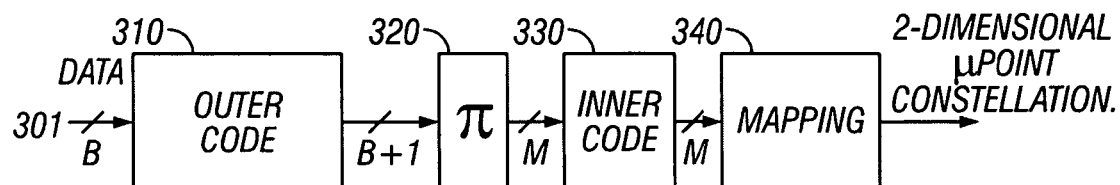


FIG. 3

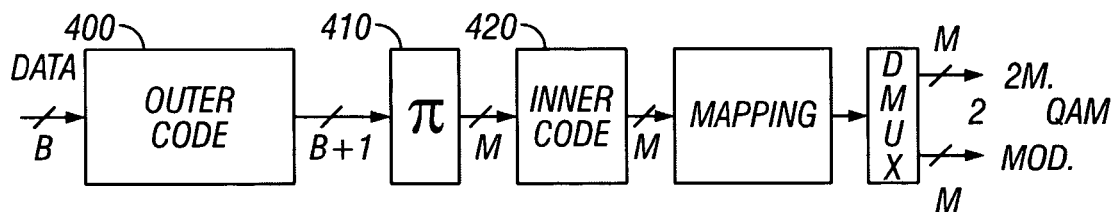


FIG. 4

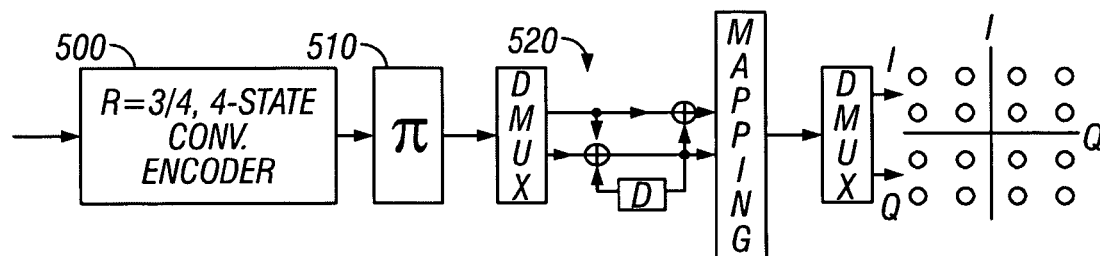


FIG. 5

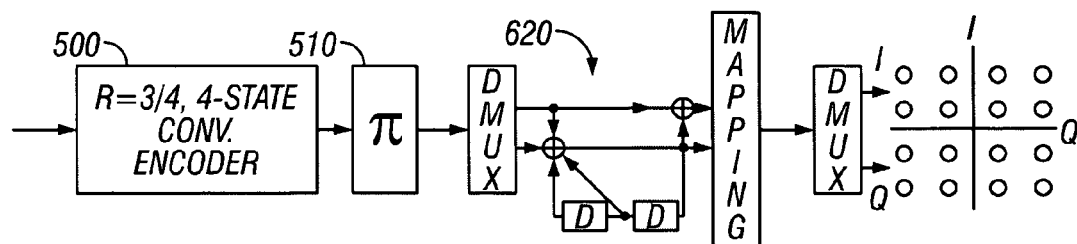


FIG. 6

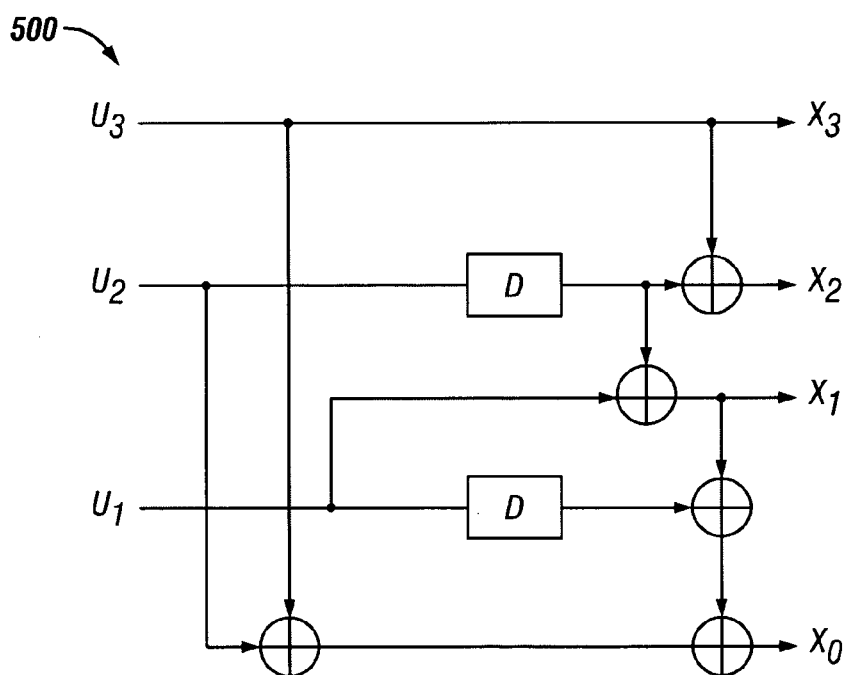


FIG. 7

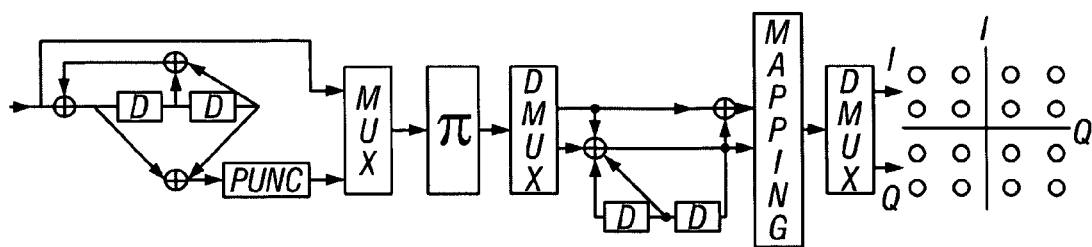


FIG. 8

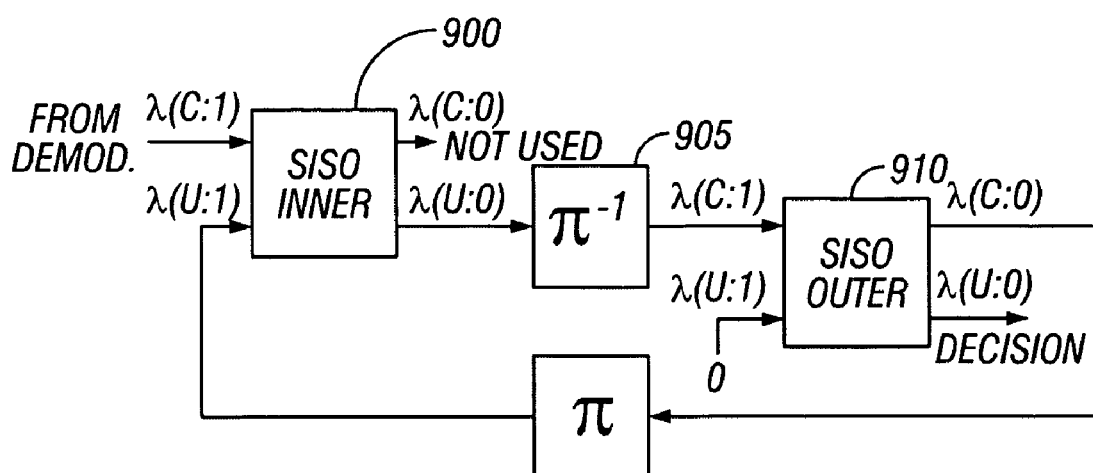


FIG. 9

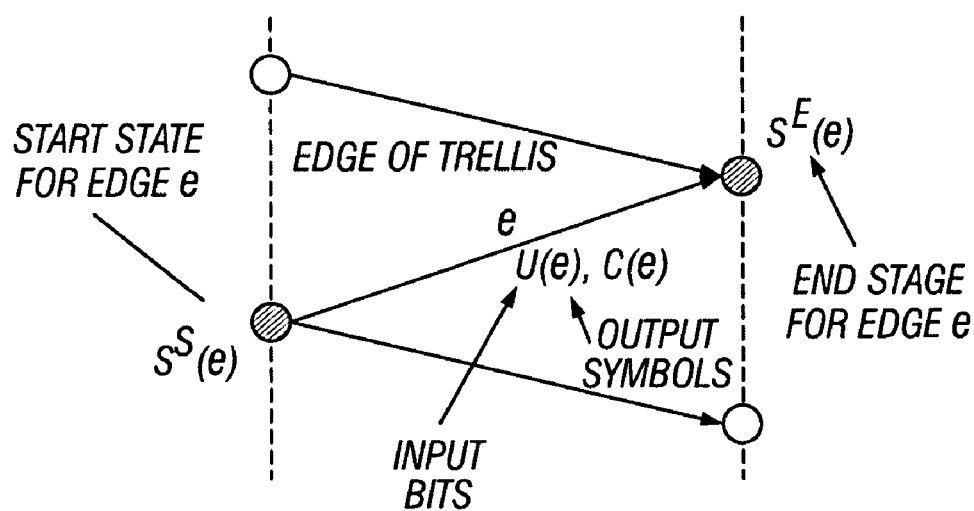


FIG. 10

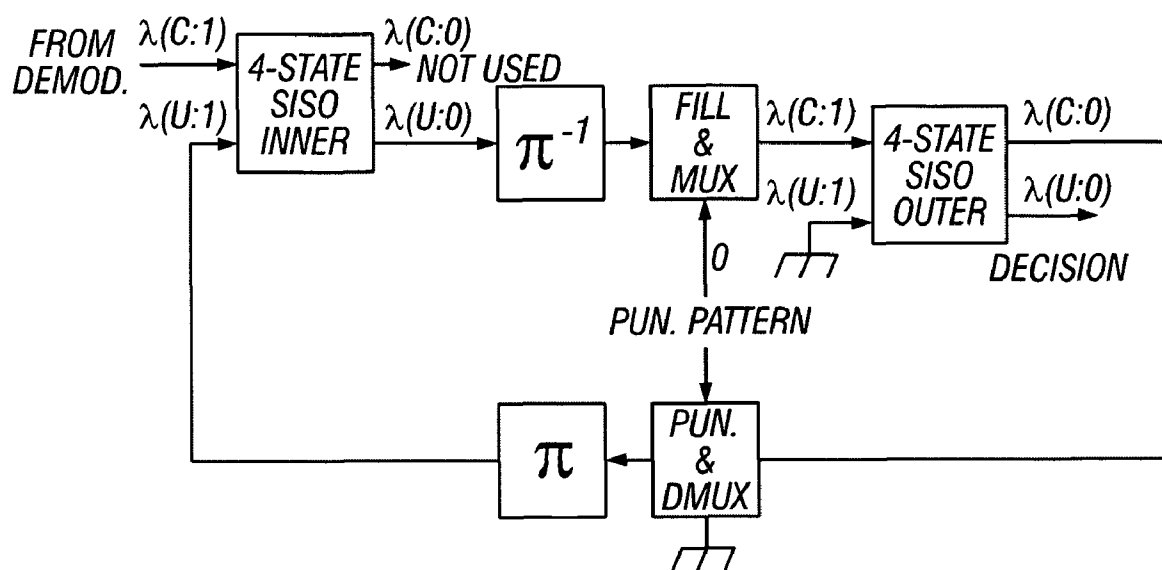


FIG. 11

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# SERIAL TURBO TRELLIS CODED MODULATION USING A SERIALY CONCATENATED CODER

## CROSS REFERENCE TO RELATED APPLICATIONS

This application claims the benefit of the U.S. Provisional Application No. 60/176,404, filed on Jan. 13, 2000.

## STATEMENT AS TO FEDERALLY-SPONSORED RESEARCH

The invention described herein was made in the performance of work under a NASA contract, and is subject to the provision of Public Law 96-517 (U.S.C. 202) in which the Contractor has elected to retain title.

## BACKGROUND

Properties of a channel affect the amount of data that can be handled by the channel. The so-called "Shannon limit" defines the theoretical limit of amount of data that a channel can carry.

Different techniques have been used to increase the data rate that can be handled by a channel. "Near Shannon Limit Error-Correcting Coding and Decoding: Turbo Codes," by Berrou et al. ICC, pp 1064-1070, (1993), described a new "turbo code" technique that has revolutionized the field of error correcting codes.

Turbo codes have sufficient randomness to allow reliable communication over the channel at a high data rate near capacity. However, they still retain sufficient structure to allow practical encoding and decoding algorithms. Still, the technique for encoding and decoding turbo codes can be relatively complex.

A standard turbo coder is shown in FIG. 1. A block of  $k$  information bits **100** is input directly to a first encoder **102**. A  $k$  bit interleaver **110** also receives the  $k$  bits and interleaves them prior to applying them to a second encoder **104**. The second encoder produces an output that has more bits than its input, that is, it is a coder with rate that is less than 1. The encoders **102**, **104** are also typically recursive convolutional coders.

Three different items are sent over the channel **150**: the original  $k$  bits **100**, first encoded bits **111**, and second encoded bits **112**.

At the decoding end, two decoders are used: a first constituent decoder **160** and a second constituent decoder **162**. Each receives both the original  $k$  bits, and one of the encoded portions **110**. Each decoder sends likelihood estimates of the decoded bits to the other decoders. The estimates are used to decode the uncoded information bits as corrupted by the noisy channel.

Turbo codes are effectively parallel concatenated codes with an encoder having two or more constituent coders joined through one or more interleavers. Input information bits feed the first encoder, are scrambled by the interleaver, and enter the second encoder. A code word is formed by a parallel concatenated code formed by the input bits to the first encoder followed by the parity check bits of both encoders.

Trellis coded modulation is described in "Channel Coding with Multilevel Phase Signaling", Ungerboeck, IEEE Trans Inf. Th. Vol. IT-25, pp 55-67, January 1982. Trellis coded modulation can produce significant coding gains in certain circumstances.

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In some situations it may be desirable to have a very low bit error rate, e.g. less than  $10^{-9}$ .

## SUMMARY

The present application combines a combination of trellis coded modulation with turbo codes, to obtain certain advantages of bandwidth and power efficiency from the trellis coded modulation, while also obtaining other advantages of the turbo codes. A specific embodiment combines serially concatenated coding for the inner coder with trellis codes on the outer coder.

## BRIEF DESCRIPTION OF THE DRAWINGS

These and other aspects of the invention will be described in detail with reference to the accompanying drawings, wherein:

FIG. 1 shows a block diagram of a prior art turbo coder;

FIG. 2 shows a block diagram of inner coder for serially concatenated trellis coded modulation using a generic mapper;

FIG. 3 shows a block diagram of an inner coder using two-dimensional  $M$  point mapping;

FIG. 4 shows a coder using a mapping system that provides trellis coded modulation for QAM;

FIG. 5 shows a trellis coded modulator which has an inner coder formed of a two state device;

FIG. 6 shows a trellis coder with a four state trellis coded modulator;

FIG. 7 shows an outer coder for use in the FIGS. 5 and 6 embodiments;

FIG. 8 shows an alternative embodiment using bit puncturing;

FIG. 9 shows a block diagram of an iterative decoder;

FIG. 10 shows a trellis diagram for the decoder; and

FIG. 11 shows a turbo coder with lower complexity.

## DETAILED DESCRIPTION

A disclosed embodiment uses serially concatenated codes with Trellis codes, to obtain low error floors and obtain the advantages of iterative coding as it is often used in a parallel concatenated code.

In a "classical" concatenated coding system, an interleaver is placed between inner and outer coders to separate bursts of errors produced by the inner encoder. In contrast, the serially concatenated coder described herein may optimize the inner and outer coders and the interleaver as a single entity thereby optimizing the whole serial structure. This has not been done in the past due to complexity and the difficulty of optimum coding.

The present application may use the technology of the uniform interleaver as described in "unveiling turbo codes: some results on parallel concatenated coding schemes", S. Benedetto, et al, IEEE TRANS of Inf Theory March 1996. The uniform interleaver allows setting criteria which optimize the component codes in order to construct more powerful serially concatenated codes with a relatively large block size.

The complexity of the coding is handled herewith using sub optimum iterative decoding methods. The concatenation of an outer convolutional code or a short block code with an inner trellis coded modulation code is called a serially concatenated TCM code. This system enables a relatively very low bit error rate.

FIG. 2 shows the basic structure of the serially concatenated trellis coded modulation scheme. The outer coder, which is a serial concatenated coder **200**, receives input data **202** having 2b bits, and produces output data **204** having 2b+1 bits. Hence, the outer coder **200** has a rate  $2b/(2b+1)$ . More generally, however, the coder should have a rate somewhat less than one. A short block code can alternatively be used as long as it has maximum free Hamming distance as the outer code.

An interleaver **210** permutes the output of the outer coder **200**. This produces interleaved data **212**. The interleaved data **212** enters an inner coding block **220** which is a recursive, convolutional inner coder having rate  $(2b+1)/(2b+2)$ . Mapper **230** then maps the 2b+2 output bits of the inner coder **220** to two symbols. Each symbol belongs to a  $2^{b+1}$  level modulation or four dimensional modulation. This system uses 2b information bits for each two modulation symbol intervals, thereby resulting in a b bit/second/Hz transmission when ideal Nyquist pulse shaping is used. In other words, this provides b bits per modulation symbol. The inner code and the mapping are jointly optimized based on maximum effective free Euclidean distance of the inner trellis coded modulation, as described above.

There are many different ways of configuring two-dimensional and multidimensional trellis coded modulators. Conventional trellis coded modulator designs may have drawbacks when used in this situation. Therefore, while the present application contemplates using conventional trellis coded modulators, it is noted that there are reasons why such conventional modulators may be less useful.

In a serial trellis coded modulator, the Euclidean distance of encoded sequences can be very large for input sequences having a Hamming distance equal to one. This may not be satisfied even if the encoder structure has feedback. Some of the input bits may remain uncoded in a conventional trellis coded modulator. These uncoded bits may select a point from among a set that has been chosen according to the encoded bits. The combination of coded and uncoded bits is then mapped to either two or higher dimensional modulation.

It has been considered by the present inventors to use conventional trellis coded modulation without parallel branches. This, however, may require that the number of states be greater than the number of transition per states. This in turn may prevent the use of simple codes with a small number of states.

Conventional trellis coded modulators also assign the input labels effectively arbitrarily. It has been thought by many that the assignment of input labels did not play an important role in coding. According to the present specified coding system, input labels are carefully selected.

Another aspect is the complexity of the code selection. The serially concatenated trellis coded modulation described with reference to FIG. 2 has a number of transitions per state of  $2^{2b+1}$ . For specific case of interest, b may equal 3. Therefore, even if the number of states is low, the number of transitions may be high. For two states, there still may be 128 transitions per state, resulting in 256 edges in the trellis section. The complexity of the decoder may depend on the number of edges per trellis section. This complexity as described above may actually interfere with high-speed operation, since the complexity of operation takes time to complete.

Another serial concatenated trellis coded modulation scheme is shown in FIG. 3. This system uses a two-dimensional constellation with M points. For purposes of explanation, we can define  $m = \log 2M$ , where M is the

number of phases. In this structure, the input data **300** is coupled to an outer coder **310** producing b+1 bits for the b input bits. Hence, the outer coder is a rate  $b/(b+1)$  binary convolutional coder. An interleaver **320** permutes the output of the outer coder. The interleaved data enters a rate  $m/m=1$  recursive convolutional inner coder. The m output bits are then mapped to one symbol along into a  $2^m$  level modulation by a mapping element **340**. This system uses b information bits per b+1/m modulation symbol interval. It effectively results in  $bm/(b+1)$  bits per modulation symbol.

The inner coder **330** and mapping **340** are jointly optimized based on maximization of the effective free Euclidean distance of the inner trellis coded modulator.

For example consider 8PSK modulation, where  $m=3$ . Then, the throughput  $r=3b/(b+1)$  is as follows: for  $b=2$ ,  $r=2$ ; for  $b=3$ ,  $r=2.25$ ; and for  $b=4$ ,  $r=2.4$ . Accordingly, a 1/2 convolutional code with puncturing can be used to obtain various throughput values, without changing the inner coder modulation.

For rectangular  $M^2$ -QAM, where  $m = \log_2 M$ , the structure may become even simpler. In this case, to achieve throughput of  $2mb/(b+1)$  bps/Hz a rate  $b/(b+1)$  outer coder and a rate  $m/m$  inner coder may be used, where the m output bits are alternatively assigned to in-phase and quadrature components of the  $M^2$ -QAM modulation.

The structure of the SCTCM encoder is shown in FIG. 4. An outer coder **400** is connected to an interleaver **410**, which drives a trellis code modulator inner coder **420**.

For example consider 16-QAM modulation, where  $m=2$ , then the throughput  $r=4b/(b+1)$  is: for  $b=1$ ,  $r=2$ ; for  $b=2$ ,  $r=2.67$ ; for  $b=3$ ,  $r=3$ ; and for  $b=4$ ,  $r=3.2$ .

For this embodiment,  $b=r=3$ . This causes the number of transitions per state of the inner TCM **420** to be reduced to 4. This results in a large reduction in complexity: 32 times lower than the previous case. Moreover, the outer coder also has a lower code rate; this code rate may be reduced from 6/7 to 3/4.

Other embodiments of this basic idea are also possible by changing the mapping. In the FIGS. 5 and 6 embodiments, the output of the inner coder is mapped to the I and Q components of 16QAM alternatively. The encoder structure of a SCTCM for 2-state inner TCM is shown in FIG. 5, which shows the rate 3/4 four state coder **500** operating as the outer coder. An interleaver **510** drives the inner coder **520**.

The encoder structure of SCTCM for 4-state inner TCM is shown in FIG. 6, with trellis coder/mapper **620**. The outer coder **500** has an optimum rate of 3/4, 4-state nonrecursive convolution code with free Hamming distance of 3.

The detailed structure of the outer encoder **500** is shown in FIG. 7. This rate 3/4, 4-state outer code has 32 edges per trellis section and produces 4 output bits. Thus the complexity per output bit is  $32/4=8$ . The complexity per input bit is 32/3.

The complexity of the outer coder may be further reduced using a rate of 1/2, 4-state systematic recursive convolutional code. This code can be punctured to rate 3/4, by puncturing only the parity bits. The minimum distance of this punctured code is 3, the same as for the optimum code. Now the code has 8 edges per trellis section and produces 2 output bits. Thus the complexity per output bit is  $8/2=4$ . Since this code is systematic there is no complexity associated with the input. The encoder structure for this low complexity SCTCM is shown in FIG. 8.



Using this low complexity scheme with 5 iterations is roughly equal to the complexity of a standard Viterbi decoder. However, this obtains a 2 db advantage over the "Pragmatic" TCM system.

It can be shown that a dominant term in the transfer function bound on bit error probability of serially concatenated TCM, employing an outer code with free (or minimum) Hamming distance  $d_f^0$ , averaged over all possible interleavers of N bits, is proportional for large N to

$$N^{-(d_f^0+1)/2} e^{-\delta^2(E_s/4N_0)}$$

Where  $\lfloor x \rfloor$  represents, the integer part of x, and

$$\delta^2 = \frac{d_f^0 d_{f,eff}^2}{2} \text{ for } d_f^0 \text{ even, and}$$

$$\delta^2 = \frac{(d_f^0 - 3) d_{f,eff}^2}{2} + (h_m^{(3)})^2, \text{ for } d_f^0 \text{ odd}$$

The parameter  $d_{f,eff}$  is the effective free Euclidean distance of the inner code,  $h_m^{(3)}$  is the minimum Euclidean distance of inner code sequences generated by input sequences with Hamming distance 3, and  $E_s/N_0$  is the M-ary symbol signal-to-noise-ratio.

The above results are valid for very large N. On the other hand, for large values of the signal-to-noise ratio  $E_s/N_0$ , the performance of SCTCM is dominated by

$$N^{-(l_m(h_m)-1)/2} e^{-h^2(E_s/4N_0)}$$

where  $h_m$  is the minimum Euclidean distance of the SCTCM scheme, and  $l_m(h_m) \geq d_f^0$ .

Based on these results, the design criterion for serially concatenated TCM for larger interleavers and very low bit error rates is to maximize the free Hamming distance of the outer code (to achieve interleaving gain), and to maximize the effective free Euclidean distance of the inner TCM code.

Let z be the binary input sequence to the inner TCM code, and x(z) be the corresponding inner TCM encoder output with M-ary symbols. The present application defines criteria for selecting the constituent inner TCM encoder:

1. The constituent inner TCM encoder may be configured for a given two or multidimensional modulation such that the minimum Euclidean distance  $d(x(z), x(z'))$  over all z, z' pairs,  $z \neq z'$  is maximized given that the Hamming distance  $d_H(z, z')=2$ . We call this minimum Euclidean distance the effective free Euclidean distance of the inner TCM code,  $d_{f,eff}$ .

2. If the free distance of outer code  $d_f^0$  is odd, then, among the selected inner TCM encoders, choose those that have the maximum Euclidean distance  $d(x(z), x(z'))$  over all z, z' pairs,  $z \neq z'$ , given that the Hamming distance  $d_H(z, z')=3$ . This value is the minimum Euclidean distance of the inner TCM code due to input Hamming distance 3, denoted by  $h_m^{(3)}$ .

3. Among the candidate encoders, select the one that has the largest minimum Euclidean distance in encoded sequences produced by input sequences with Hamming distance  $d_f^0$ . This minimum Euclidean distance of the SCTCM is called  $h_m$ .

It has been found by the inventors that that sequences with Hamming distances of 2 or 3 at the input of the TCM encoder are still important, even if the free Hamming distance  $d_f^0$  of the outer code is larger than 2 or even 3. This is because the interleaving gain at low signal to noise ratios may depend on the number of error events that a pair of input

sequences generate in the trellis of the inner code. For a given input Hamming distance, a larger number of error events may create a smaller interleaving gain. For example, if the input Hamming distance between sequences to the inner TCM is 4, the largest number of error events that produce small output Euclidean distances is 2 (two events with an input Hamming distance of 2 each).

As described above, the present embodiments also use mapping of output labels for TCM. As soon as the input labels and output signals are assigned to the edges of a trellis, a complete description of the TCM code is obtained. The selection of the mapping (output labels) does not change the trellis code. However, it influences the encoder circuit required to implement the TCM scheme. A convenient mapping should be selected to simplify the encoder circuit and, if possible, to yield a linear circuit that can be implemented with exclusive Ors. The set partitioning of the constellation and the assignment of constellation points to trellis edges, and the successive assignments of input labels to the edges may be important. Ungerboeck proposed a mapping called "Mapping by set partitioning", leading to the "natural mapping". This mapping for two-dimensional modulation may be useful if one selects the TCM scheme by searching among all encoder circuits that maximize the minimum Euclidean distance.

The "inner" trellis code modulator can be configured as follows:

The well known set partitioning techniques for signal sets may be used.

The input label assignment is based on the codewords of the parity check code (m, m-1, 2) and the set partitioning, to maximize the quantities described in the equations above. The minimum Hamming distance between input labels for parallel transitions will be equal to 2. The assignment of codewords of the parity check code as input labels to the two-dimensional signal points is not arbitrary.

A sufficient condition to have very large output Euclidean distances for input sequences with Hamming distance 1 is that all input labels to each state be distinct.

A pair of input labels and two-dimensional signal points are assigned to the edges of a trellis diagram based on the design criteria described above.

#### EXAMPLE 1

##### Set Partitioning of 8PSK and Input Labels Assignment

Let the eight phases of 8PSK be denoted by {0, 1, 2, 3, 4, 5, 6, 7}. Here m=3. Consider the 8PSK signal set A={0, 2, 4, 6}, and set B={1, 3, 5, 7}. For unit radius 8PSK constellation, the minimum intra-set square Euclidean distance for each set is 2. The minimum inter-set square Euclidean distances is 0.586.

Select the input label set  $L_0$  as codewords of the (3, 2, 2) parity check code, i.e.  $L_0 = \{(000), (011), (101), (110)\}$ , next generate input label  $L_1 = L_0 + (001)$ , i.e.,  $L_1 = \{(001), (010), (100), (111)\}$ . Consider a 2-state trellis. Assign the input-output pair ( $L_0$ , A) to four edges from state 0 to state 0. Assign the input-output pair ( $L_1$ , B) to four edges from state 0 to state 1. Next assign the input-output pair ( $L_2$ , A) to four edges from the state 1 to state 0, and assign the input-output pair ( $L_3$ , B) to four edges from—state 1 to state 1.  $L_2$  has the same elements as in  $L_1$  but with different order, and  $L_3$  has the same elements as in  $L_0$  again with different order. In order to maximize the minimum Euclidean distance due to

the input sequences with Hamming distance 2, we have to find the right permutation within each set. In this case it turns out that using the complement operation suffices. Therefore define input label  $L_2$  as the complement of the elements of  $L_0$  without changing the order, i.e.,  $L_2 = [(111), (100), (010), (001)]$ . Finally  $L_3$  is generated in the same way, as the complement of elements in  $L_1$ , i.e.  $L_3 = [(110), (101), (011), (000)]$ .

Such assignment guarantees that the squared effective free Euclidean distance of trellis code is 2, where the minimum squared Euclidean distance of the code is 0.586.

Having determined the code by its input labels and two-dimensional output signals, the encoder structure can then be obtained by selecting any appropriate labels (output labels) for the two-dimensional output signals. The following output mapping may be used:  $\{(000), (001), (010), (011), (110), (111), (100), (101)\}$ , mapped to phases  $[0, 1, 2, 3, 4, 5, 6, 7]$ , which is called "reordered mapping". For this 2-state inner code,  $d_{f,eff}^2=2$ , and  $h_m^{(3)}=\infty$ , and  $h_m^2=0.586$ . The outer code for this example can be selected as a 4-state, rate 2/3, convolutional code with  $d_p^0=3$  (this is a recursive systematic rate 1/2 convolutional code where the parity bits are punctured). Since  $h_m^{(3)}=\infty$  then  $d_f^0$  is increased effectively to 4. This method of design was used to obtain the encoders in the previous examples for 16QAM.

A decoder is described herein. This decoder can be a Bit-by-Bit Iterative Decoder. The iterative decoder for serially concatenated trellis coded modulation uses a generalized Log-APP (a-posteriori probability) decoder module with four ports, called SISO APP module or simply SISO. The block diagram of the iterative decoder for serial concatenated TCM is shown in FIG. 9. The device has a SISO inner decoder **900** coupled to a deinterleaver **905**, an outer decoder **910**. Feedback is passed through an interleaver **920** back to the inner decoder.

The decoding techniques may be used for the inner TCM code and outer convolutional code, using the trellis section shown in FIG. 10. Consider an inner TCM code with  $p_1$  input bits and  $q_1$  nonbinary complex output symbols with normalized unit power, and an outer code with  $p_2$  input bits and  $q_2$  binary outputs  $\{0, 1\}$ . Let  $U_k(e)$  represent  $u_{k,i}(e)$ ;  $i=1, 2, \dots, p_m$  the input bits on a trellis edge at time  $k$  ( $m=1$  for the inner TCM, and  $m=2$  for the outer code), and let  $c_{k,i}(e)$  represents  $c_{k,i}(e)$ ;  $i=1, 2, \dots, q_m$  the output symbols ( $m=1$  for the inner TCM, with nonbinary complex symbols, and  $m=2$  for the outer code with binary  $\{0, 1\}$  symbols).

Define the reliability of a bit  $Z$  taking values  $\{0, 1\}$  at time  $k$  as

$$\lambda_k[z; \dots] \Delta \log \frac{P_k[Z=1; \cdot]}{P_k[Z=0; \cdot]}$$

The second argument in the brackets, shown as a dot, may represent I, the input, or O, the output, to the SISO. We use the following identity

$$a = \log \left[ \sum_{i=1}^L e^{a_i} \right] = \max_i \{a_1 \dots, a_L\} \Delta \max_i \{a_i\}$$

where  $\delta(a_1, \dots, a_L)$  is the correction term which can be computed using a look-up table.

The "max\*" operation is a maximization (compare/select) plus a correction term (lookup table). Small degradations occur if the "max\*" operation is replaced by "max". The received complex samples  $\{y_{k,i}\}$  at the output of the receiver matched filter are normalized such that additive complex noise samples have unit variance per dimension.

SISO can be used for the Inner TCM.

The forward and the backward recursions are:

$$\alpha_k(s) =$$

$$\max_{e: s^E(e)=s} * \left\{ \alpha_{k-1}[S^S(e)] + \sum_{i=1}^{p_1} u_{k,i}(e) \lambda_k[U_{k,i}; I] + \sum_{i=1}^{q_1} \bar{\lambda}_k[c_{k,i}(e); I] \right\} + h_{\alpha k}$$

$$\beta_k(s) = \max_{e: s^S(e)=s} * \left\{ \beta_{k+1}[S^E(e)] + \right.$$

$$\left. \sum_{i=1}^{p_1} u_{k+1,i}(e) \lambda_{k+1}[U_{k+1,i}; I] + \sum_{i=1}^{q_1} \bar{\lambda}_{k+1}[c_{k+1,i}(e); I] \right\} + h_{\beta k}$$

for all states  $s$ , and  $k=1, \dots, (n-1)$ , where  $n$  represents the total number of trellis steps from the initial state to the final state.

The extrinsic bit information for  $U_{k,j}$ ;  $j=1, 2, \dots, p_1$  can be obtained from:

$$\lambda_k(U_{k,j}; O) = \max_{e: u_{k,j}(e)=1} * \left\{ \alpha_{k-1}[S^S(e)] + \sum_{i=1, i \neq j}^{p_1} u_{k,i}(e) \lambda_k[U_{k,i}; I] + \right.$$

$$\left. \sum_{i=1}^{q_1} \bar{\lambda}_k[c_{k,i}(e); I] + \beta_k[S^E(e)] \right\} - \max_{e: u_{k,j}(e)=0} *$$

$$\left\{ \alpha_{k-1}[S^S(e)] + \sum_{i=1, i \neq j}^{p_1} u_{k,i}(e) \lambda_k[U_{k,i}; I] + \sum_{i=1}^{q_1} \bar{\lambda}_k[c_{k,i}(e); I] + \beta_k[S^E(e)] \right\}$$

$$\text{where } \bar{\lambda}_k[c_{k,i}(e); I] = - \left| y_{k,i} - \sqrt{\frac{2E_s}{N_0}} c_{k,i}(e) \right|^2 / 2.$$

We assume the initial and the final states of the inner encoder (as well as the outer encoder) are the all zero state. Forward recursions start with initial values,  $\alpha_0(s)=0$ , if  $s=0$  (initial zero state) and  $\alpha_0(s)=-\infty$ , if  $s \neq 0$ . Backward recursions start with  $\beta_n(s)=0$ , if  $s=0$  (final zero state) and  $\beta_n(s)=-\infty$ , if  $s \neq 0$ . The  $h_{\alpha k}$  and  $h_{\beta k}$  are normalization constants. Which, in the hardware implementation of the SISO, are used to prevent buffer overflow. These operations are similar to the Viterbi algorithm used in the forward and backward directions, except for a correction term that is added when compare-select operations are performed. At the first iteration, all  $\lambda_k[U_{k,i}; I]$  are zero. After the first iteration, the inner SISO accepts the extrinsics from the outer SISO, through the interlayer  $\pi$ , as reliabilities of input bits of TCM encoder, and the external observations from the channel. The inner SISO uses the input reliabilities and observations for the calculation of new extrinsics  $\lambda_k(U_{k,j}; O)$  for the input bits. These are then provided to the outer SISO module, through the deinterleaver  $\pi^{-1}$ . The forward and the backward recursions for SISO are:

$$a_k(s) = \max_{e: s^E(e)=s} * \left\{ a_{k-1}[s^S(e)] + \sum_{i=1}^{q_2} c_{k,i}(e) \lambda_k[C_{k,i}; I] \right\} + h_{a_k}$$

$$\beta_k(s) = \max_{e: s^E(e)=s} * \left\{ \beta_{k+1}[s^E(e)] + \sum_{i=1}^{q_2} c_{k+1,i}(e) \lambda_k[C_{k+1,i}; I] \right\} + h_{\beta_k}$$

The extrinsic information for  $C_{k,j}$ ;  $j=1, 2, \dots, q_2$ , can be obtained from:

$$\lambda_k(C_{k,j}; O) = \max_{e: c_{k,j}(e)=1} * \left\{ a_{k-1}[s^S(e)] + \sum_{\substack{i=1 \\ i \neq j}}^{q_2} c_{k,i}(e) \lambda_k[C_{k,i}; I] + \beta_k[s^E(e)] \right\} -$$

$$\max_{e: c_{k,j}(e)=0} * \left\{ a_{k-1}[s^S(e)] + \sum_{\substack{i=1 \\ i \neq j}}^{q_2} c_{k,i}(e) \lambda_k[C_{k,i}; I] + \beta_k[s^E(e)] \right\}$$

with initial values,  $a_0(s)=0$ , if  $s=0$  and  $a_0(s)=-\infty$ , if  $s \neq 0$  and  $\beta_n(s)=0$ , if  $s=0$  and  $\beta_n(s)=-\infty$ , if  $s \neq 0$ , where  $h_{a_k}$  and  $h_{\beta_k}$  are normalization constants which, in the hardware implementation of the SISO, are used to prevent the buffer overflow.

The final decision is obtained from the bit reliability computation of  $U_{k,j}$ ;  $j=1, 2, \dots, p_2$ , passing through a hard limiter, as

$$\lambda_k(U_{k,j}; O) = \max_{e: u_{k,j}(e)=1} * \left\{ a_{k-1}[s^S(e)] + \sum_{i=1}^{q_2} c_{k,i}(e) \lambda_k[C_{k,i}; I] + \beta_k[s^E(e)] \right\} -$$

$$\max_{e: u_{k,j}(e)=0} * \left\{ a_{k-1}[s^S(e)] + \sum_{i=1}^{q_2} c_{k,i}(e) \lambda_k[C_{k,i}; I] + \beta_k[s^E(e)] \right\}$$

The outer SISO accepts the extrinsics from the inner SISO as input reliabilities of coded bits of the outer encoder. For the outer SISO there is no external observation from the channel. The outer SISO uses the input reliabilities for calculation of new extrinsics  $\lambda_k(C_{k,j}; O)$  for coded bits. These are then provided to the inner SISO module.

The structure of iterative decoder for punctured outer code is shown in FIG. 11.

Other embodiments are within the disclosed invention.

What is claimed is:

1. An encoding system configured to output symbols, comprising:

an outer coder, configured to receive data to be coded, and producing outer coded data as an output;

an interleaver, coupled to receive said outer coded data as an input, and to permute said outer coded data to produce interleaved data;

a recursive rate m/m inner coder, receiving said interleaved data, and coding m bits of said interleaved data to produce m bits of inner coded output data according to an inner code, wherein said recursive rate m/m inner coder has a structure which facilitates iterative decoding of said symbols at a decoder;

a mapping element, receiving an output of said recursive rate m/m inner coder, and assigning said m output bits to a first of said symbols according to a mapping;

wherein said recursive rate m/m inner coder and said mapping element are selected to maximize the effective free Euclidean distance of a trellis coded modulator formed from said recursive rate m/m inner coder and said mapping element.

2. The system of claim 1, wherein said first symbol is a point in a QAM constellation.

3. The system of claim 1, further comprising a demultiplexer, wherein the demultiplexer is configured to:

receive successive ones of said symbols including said first symbol from the mapping element; and

provide the successive symbols to I and Q channels of a quadrature amplitude modulator in an alternating fashion.

4. A serially concatenated trellis coded modulation (SCTCM) encoder configured to output symbols, comprising:

an outer coder configured to receive input data, code said input data according to a first code, and produce outer coded data;

an interleaver configured to receive said outer coded data, permute said outer coded data, and produce interleaved data;

a trellis coded modulator (TCM), including:

a recursive rate m/m coder configured to receive m bits of said interleaved data, code said m bits of interleaved data according to a second code, and produce m bits of inner coded data, wherein said recursive rate m/m coder has a structure which facilitates iterative decoding of said symbols at a decoder;

a mapping element configured to map said m bits of inner coded data to a first of said symbols according to a mapping;

wherein said recursive rate m/m coder and said mapping element maximize the effective free Euclidean distance of the TCM.

5. The SCTCM encoder of claim 4, wherein the outer coder is a rate b/(b+1) coder, wherein b is a positive integer.

6. The SCTCM encoder of claim 4, further comprising a demultiplexer coupled to said TCM, wherein said demultiplexer is configured to receive said symbols from said TCM and alternately assign said symbols to an in-phase (I) or quadrature (Q) component of a QAM modulation.

7. The SCTCM encoder of claim 4, wherein said symbols correspond to an MPSK modulation.

8. The SCTCM encoder of claim 4, wherein said symbols correspond to an  $M^2$  QAM modulation.

9. The SCTCM encoder of claim 8, wherein said symbols correspond to a 16 QAM modulation.

10. A serially concatenated trellis coded modulation (SCTCM) encoder, comprising:

an outer coder configured to receive input data, code said input data according to a first code, and produce outer coded data;

an interleaver configured to receive said outer coded data, permute said outer coded data, and produce interleaved data;

a trellis coded modulator (TCM) configured to receive said interleaved data from said interleaver, and wherein said TCM is configured to receive a bit sequence z of interleaved data and produce a corresponding sequence of output symbols x(z), said TCM including:

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a recursive rate-1 inner coder configured to receive said interleaved data, code said interleaved data according to a second code, and produce inner coded data, wherein said recursive rate-1 inner coder has a structure which facilitates iterative decoding of said symbols at a decoder;

a mapper configured to receive bits of said inner coded data and map said bits of said inner coded data according to a mapping to produce a first of said symbols;

wherein said second code and said mapping are selected to maximize the minimum Euclidean distance  $d(x(z), x(z'))$ , wherein the minimum Euclidean distance to be maximized is determined over all pairs  $z, z'$ , wherein  $z'$  is a bit sequence for which  $z \neq z'$  and the Hamming distance  $d_H(z, z')$  is 2.

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11. The SCTCM encoder of claim 10, wherein the outer coder is a rate  $b/(b+1)$  coder, wherein  $b$  is a positive integer.
12. The SCTCM encoder of claim 10, wherein the first code is a convolutional code.
13. The SCTCM encoder of claim 10, wherein the first code is a block code.
14. The SCTCM encoder of claim 10, wherein said symbols correspond to an MPSK modulation.
15. The SCTCM encoder of claim 10, wherein said symbols correspond to a  $M^2$  QAM modulation.
16. The SCTCM encoder of claim 15, wherein said symbols correspond to a 16 QAM modulation.

\* \* \* \* \*